

J213 Paint the Wall

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香港電腦奧林匹克競賽
Hong Kong Olympiad in Informatics

Problem

Given a $R \times C$ grid, initially all cells are white, paint K cells to black. Maximize pairs of adjacent cells with different color.

3 3 3

010

100

010

For all cases:

$$1 \leq R, C \leq 100$$

$$0 \leq N \leq R \times C$$

	Points	Constraints
1	13	$R = 1$
2	18	$R = 2$ $2 \leq C \leq 100$
3	9	$R = C = 3$
4	8	$R \times C$ is even $N = \frac{R \times C}{2}$
5	29	$R \times C$ is even
6	23	No additional constraints



Stats

First solve: cwong 1:14

7 contestants had scored 100

Mean: 22

Subtask 3

$$R = C = 3$$

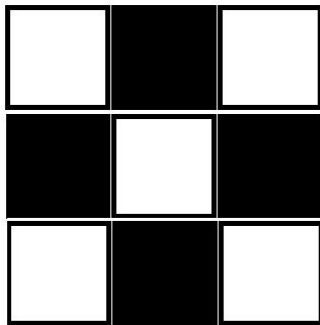
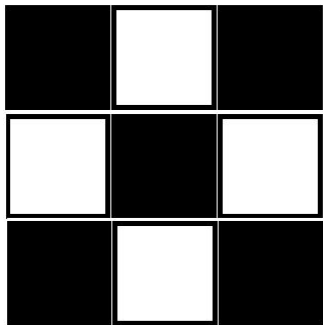
You can solve $K=1..9$ on paper and hardcode it, or writing a brute force algorithm to generate all possible colourings and find the optimal one.

Time complexity: $O(1)$ or $O(2^{R * C})$

Main Observation 1

Notice that when $K > R * C / 2$, the problem can be transformed to, initially all cells are '1', we are changing $R * C - K$ cells to '0'. So, we can solve the original problem with $K = R * C - K$ and flip the cell color at last.

From now on, we assume that $K \leq R * C / 2$.



Subtask 1

$$R = 1$$

Intuitively, we know when C is odd, we should choose cell 2, 4, ..., $C - 1$. (row 1)

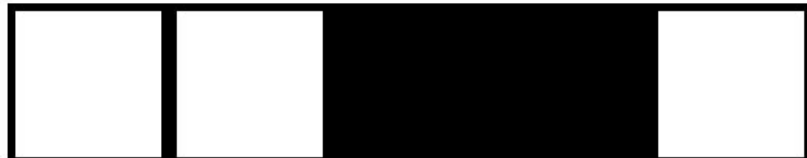
And when C is even, we can choose 1, 3, ..., $C - 1$ or 2, 4, ..., C .

Why the parity of C matters? Think about $C = 3$, we have to choose the middle cell since only it has two neighbours. In general, we don't really want to choose cell 1 or C unless we have no choice.



Main Observation 2

When $K \leq R * C / 2$, in optimal answer, we will never paint two adjacent cells with '1'. We can always construct such answer (choose odd or even columns).



Subtask 1

Since in both cases (C is even or odd), start choosing from 2 is optimal.

So we will paint cell 2, 4, 6, ... and stop when we have painted K cells.

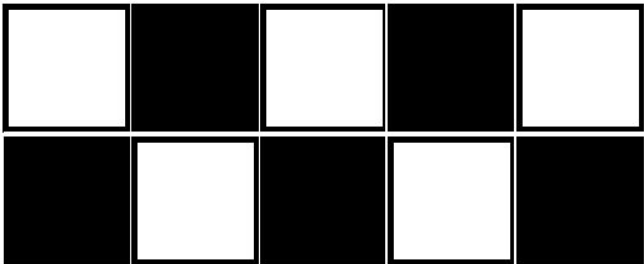
Time complexity: $O(C)$

Subtask 2

$R = 2, C \leq 100$

We can extend main observation 2, and it also works when $R = 2$.

When $K = R * C / 2$, we know our answer will be:



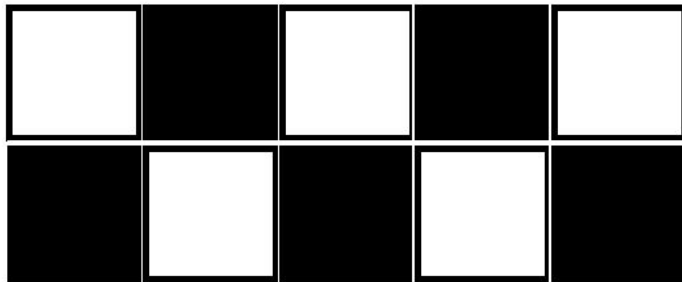
What if $K < R * C / 2$?

Main Observation 3

We can pick the cells greedily.

When $R = 2$ and $C = 5$, we have 5 choices to paint.

Notice that we can consider these choices independently and it wouldn't affect others, as we would never paint both adjacent cells with '1'.

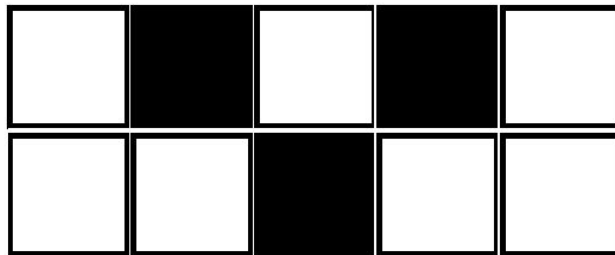


Subtask 2

We want to greedily paint cells that have more adjacent neighbours (cells that aren't located in column 1 or C).

For example, when $R = 2$, $C = 5$ and $K = 3$, the solution below is one of the optimal solutions.

Time complexity: $O(R * C)$



Subtask 4

$R * C$ is even and $K = R * C / 2$.

With our intuition or the observation we have, main observation 2 and indeed it works in general case ($R > 2$), we can notice that we will be painting the grid like this:

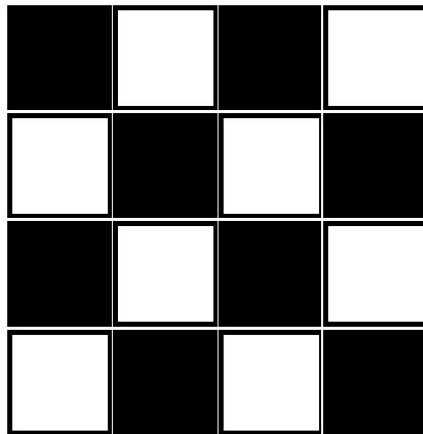
(1, 1), (1, 3), ...

(2, 2), (2, 4), ...

(3, 1), (3, 3), ...

...

Time Complexity: $O(R * C)$



Subtask 5

$R * C$ is even

We can combine our idea in subtask 2 and 4. We are picking K non-adjacent cells to paint them as '1' and we are picking them greedily by their number of adjacent neighbours. So we are picking K cells from here:

(1, 1), (1, 3), ...

(2, 2), (2, 4), ...

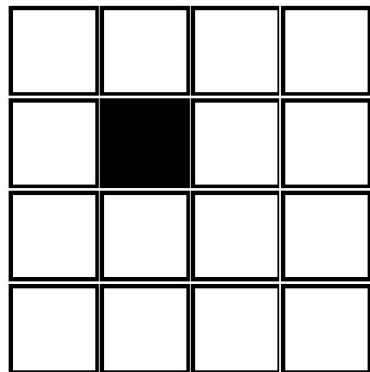
(3, 1), (3, 3), ...

...

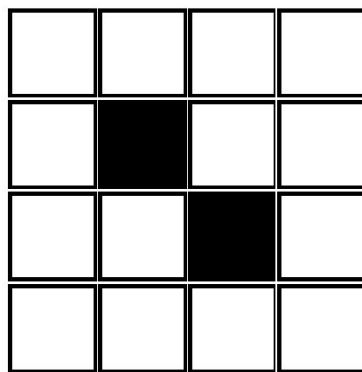


Subtask 5

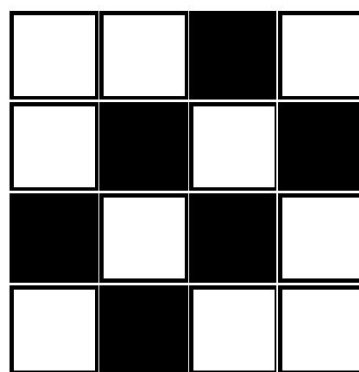
$K = 1$



$K = 2$



$K = 6$



Time complexity: $O(R * C)$

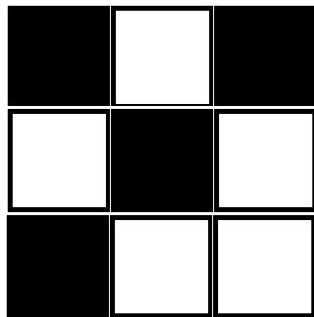
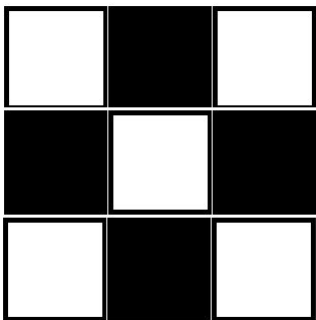
Subtask 6

No additional constraints

Why doesn't subtask 5's idea work in general?

When $R = C = 3$, $K = 4$:

it is better to paint it in the way of the left one than the right one.



Main Observation 4

We have two (and only two) different choosing mechanisms for non-adjacent cells:

1. choose cell (i, j) where $i + j = 0 \pmod{2}$
2. choose cell (i, j) where $i + j = 1 \pmod{2}$

When $R * C$ is even, two methods are the same. (Imagine $R = 4$ and $C = 4$)

When $R * C$ is odd, one might yield a better result.



Subtask 6

We try both methods and pick the one with larger result.

Time complexity: $O(R * C)$