

J211 - Paint Shop

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Background

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Task

7 colors of paints: cyan, magenta, yellow, black, red, green, and blue

Each paint price is an integer price per litre.

Obtain at least **N** litres of paint of a given color by buying and/or mixing the paints using the following formulas:

1L cyan + 1L magenta → 2L blue

1L magenta + 1L yellow → 2L red

1L cyan + 1L yellow → 2L green

1L cyan + 1L magenta + 1L yellow → 3L black

1L red + 1L green + 1L blue → 3L black

1L cyan + 2L red → 3L black

1L magenta + 2L green → 3L black

1L yellow + 2L blue → 3L black

Sample 1

	cyan	magenta	yellow	black	red	green	blue
Price	10	10	10	15	15	15	15

Required: 7 litres of red

3L magenta + 3L yellow → 6L red
1L red

$$\begin{aligned} \text{Answer} &= 3 * 10 + 3 * 10 + 15 \\ &= 75 \end{aligned}$$

1L cyan + 1L magenta → 2L blue

1L magenta + 1L yellow → 2L red

1L cyan + 1L yellow → 2L green

1L cyan + 1L magenta + 1L yellow → 3L black

1L red + 1L green + 1L blue → 3L black

1L cyan + 2L red → 3L black

1L magenta + 2L green → 3L black

1L yellow + 2L blue → 3L black



Sample 2

	cyan	magenta	yellow	black	red	green	blue
Price	10	10	10	25	25	25	25

Required: 7 litres of red

4L magenta + 4L yellow \rightarrow 8L red

$$\begin{aligned} \text{Answer} &= 4 * 10 + 4 * 10 \\ &= 80 \end{aligned}$$

1L cyan + 1L magenta \rightarrow 2L blue

1L magenta + 1L yellow \rightarrow 2L red

1L cyan + 1L yellow \rightarrow 2L green

1L cyan + 1L magenta + 1L yellow \rightarrow 3L black

1L red + 1L green + 1L blue \rightarrow 3L black

1L cyan + 2L red \rightarrow 3L black

1L magenta + 2L green \rightarrow 3L black

1L yellow + 2L blue \rightarrow 3L black

Sample 3

	cyan	magenta	yellow	black	red	green	blue
Price	1	2	3	4	3	2	1

Required: **3** litres of black

1L yellow + 2L blue → 3L black

$$\begin{aligned} \text{Answer} &= 3 + 2 * 1 \\ &= 5 \end{aligned}$$

1L cyan + 1L magenta → 2L blue

1L magenta + 1L yellow → 2L red

1L cyan + 1L yellow → 2L green

1L cyan + 1L magenta + 1L yellow → 3L black

1L red + 1L green + 1L blue → 3L black

1L cyan + 2L red → 3L black

1L magenta + 2L green → 3L black

1L yellow + 2L blue → 3L black

Subtasks

For all cases:

$$1 \leq N \leq 100000$$

$$1 \leq \text{price of any color per litre} \leq 100$$

	Points	Constraints			
1	11	$1 \leq N \leq 10$ S is always cyan, magenta or yellow.	4	19	$1 \leq N \leq 10$ N is a multiple of 3. S is always black.
2	14	$1 \leq N \leq 10$ N is an even number. S is always red.	5	27	$1 \leq N \leq 10$ S is always black.
3	23	$1 \leq N \leq 10$ S is always red, green or blue.	6	6	No additional constraints



Statistics

Task	Attempts	Max	Mean	Std Dev	Subtasks					
J211 - Paint Shop	75	100	76.84	33.51	11: 69	14: 68	23: 61	19: 57	27: 48	6: 45

First solved by **ryanjz2024** at **0:14:39**

Last solved by **dbsic** at **2:29:52** (8 seconds before the contest ends)

Easiest task among all three junior tasks

More **Accepted**, More Happiness ^_^



Solutions

Subtask 1

$$1 \leq N \leq 10$$

The color is always cyan, magenta or blue

- We can only obtain these 3 colors by buying them from the shop
- Answer = price of the color * **N**

Subtask 2

$$1 \leq N \leq 10$$

N is an even number

The color is always red

- Option 1: buy N litres of red directly
- Option 2: buy $N / 2$ litres magenta + $N / 2$ litres yellow

- Answer = min(Option 1, Option 2)

1L cyan + 1L magenta → 2L blue

1L magenta + 1L yellow → 2L red

1L cyan + 1L yellow → 2L green

Subtask 3

$$1 \leq N \leq 10$$

The color is always red, green or blue

Using red as example:

- Option 1: buy N litres of red directly
- Option 2: buy $\text{ceil}(N / 2)$ litres of magenta + $\text{ceil}(N / 2)$ litres of yellow

- Enough?

1L cyan + 1L magenta → 2L blue

1L magenta + 1L yellow → 2L red

1L cyan + 1L yellow → 2L green



Subtask 3

- Of course not, sample 1 already reminds you $\wedge _ \wedge$
- For **N** is even, same as subtask 2
- For **N** is odd, we consider the followings:
 - First **N - 1** litres: $\min((\mathbf{N} - \mathbf{1}) * \mathbf{red}, (\mathbf{N} - \mathbf{1}) / 2 * (\mathbf{magenta} + \mathbf{yellow}))$
 - Last **1** litre: $\min(\mathbf{red}, \mathbf{magenta} + \mathbf{yellow})$
- Answer = answer for **N - 1** litres + answer for **1** litre
- Same way to deal with green and blue



Subtask 4

$1 \leq N \leq 10$

N is a multiple of 3

The color is always black

- How many options this time?

1L cyan + 1L magenta → 2L blue

1L magenta + 1L yellow → 2L red

1L cyan + 1L yellow → 2L green

1L cyan + 1L magenta + 1L yellow → 3L black

1L red + 1L green + 1L blue → 3L black

1L cyan + 2L red → 3L black

1L magenta + 2L green → 3L black

1L yellow + 2L blue → 3L black



Subtask 4

- First, undoubtedly, the 5 given formulas
- But then seems more complicated formulas exist!?

1L cyan + 1L magenta → 2L blue

1L magenta + 1L yellow → 2L red

1L cyan + 1L yellow → 2L green

- 1L magenta + 1L yellow → 2L red
- Take 1L red from above, 1L red + 1L green + 1L blue → 3L black
- ...

1L cyan + 1L magenta + 1L yellow → 3L black

1L red + 1L green + 1L blue → 3L black

1L cyan + 2L red → 3L black

1L magenta + 2L green → 3L black

1L yellow + 2L blue → 3L black



Subtask 4

- In fact, the given formulas are sufficient to get optimal answer
- Do you know why?

Example

Let's prove that **magenta + yellow + green + blue** isn't optimal as an example

Denote **m = magenta**, **y = yellow**, **g = green**, **b = blue**

- Suppose $g < b$, we have $(m + y + g + b) > (m + y + g + g) > (m + g + g)$
- We can see that $(m + g + g)$ is a better formula
- Similar proof can be done for $g \geq b$

1L cyan + 1L magenta + 1L yellow → 3L black

1L red + 1L green + 1L blue → 3L black

1L cyan + 2L red → 3L black

1L magenta + 2L green → 3L black

1L yellow + 2L blue → 3L black

Remaining proof is left as exercise

Subtask 4

- Knowing that the given formulas are sufficient to get optimal answer, we just need to consider all of them & consider buying black directly
 - Option 1: $N / 3 * (\text{cyan} + \text{magenta} + \text{yellow})$
 - Option 2: $N / 3 * (\text{red} + \text{green} + \text{blue})$
 - Option 3: $N / 3 * (\text{cyan} + 2 * \text{red})$
 - Option 4: $N / 3 * (\text{magenta} + 2 * \text{green})$
 - Option 5: $N / 3 * (\text{yellow} + 2 * \text{blue})$
 - Option 6: $N * \text{black}$

1L cyan + 1L magenta + 1L yellow → 3L black

1L red + 1L green + 1L blue → 3L black

1L cyan + 2L red → 3L black

1L magenta + 2L green → 3L black

1L yellow + 2L blue → 3L black

- Answer = min(Option 1, Option 2, ... , Option 6)

Subtask 5

$$1 \leq N \leq 10$$

The color is always black

- **N** is not a multiple of **3** anymore
- What do we need to extra handle this time?



Subtask 5

- For the remainder of $N / 3 = 0$, same as subtask 4
- For the remainder of $N / 3 = 1$, what should we do?

Subtask 5

- Using similar idea from subtask 3, we consider two parts:
 - First ($N - 1$) litres: same as subtask 4
 - Last **1** litre: $\min(\text{Option 1, Option 2, ... , Option 6})$
 - Option 1: **cyan + magenta + yellow**
 - Option 2: **red + green + blue**
 - Option 3: **cyan + 2 * red**
 - Option 4: **magenta + 2 * green**
 - Option 5: **yellow + 2 * blue**
 - Option 6: **black**

- Answer = answer for $N - 1$ litres + answer for **1** litre



Subtask 5

- The same idea can be applied when the remainder of $N / 3 = 2$:
 - First ($N - 2$) litres: same as subtask 4
 - Last 2 litre: min(Option 1, Option 2, ... , Option 6)
 - Option 1: **cyan + magenta + yellow**
 - Option 2: **red + green + blue**
 - Option 3: **cyan + 2 * red**
 - Option 4: **magenta + 2 * green**
 - Option 5: **yellow + 2 * blue**
 - Option 6: **2 * black**

- Answer = answer for $N - 2$ litres + answer for 2 litre



Subtask 6

No additional constraints

- Combine subtask 1 to 5 to get the last 6 points!

Some Careless Mistakes

Pascal:

- Using `integer` but not `longint`, losing last 6 points :(
- `read(N, Col) / readln(N, Col)` leads to extra space
 - `Col = ' cyan'`

C++:

- Using `double` to output, resulting in wrong format
 - `2.9232e+06`

Implementation Tricks

For finding minimum, instead of writing brunches of **if**'s, we can make good use of **min** function and **arrays**

Example of subtask 5 when remainder of $N / 3 = 1$:

```
opt[1] = c + m + y;
opt[2] = r + g + b;
opt[3] = c + 2 * r;
opt[4] = m + 2 * g;
opt[5] = y + 2 * b;

ans[1] = (N - 1) * k;
for (int i = 1; i <= 5; i++)
    ans[1] = min(ans[1], (N - 1) / 3 * opt[i]);

ans[2] = k;
for (int i = 1; i <= 5; i++)
    ans[2] = min(ans[2], opt[i]);
```

