Snakes and Snakes

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Statistics

Score Distribution



Statistics

- Attempts : 78
- Max : 100
- Mean : 28.435
- Std Dev : 29.331

- First Solved by Jamie Choi at 0:25
- Highest mean among senior problems

Problem Statment

- Given a board consists of N+1 squares and M "Snakes"
- Each time you can move forward 1 K squares
- Find the minimum number of moves needed to move from square 0 to square N
- And determine whether you can do it

For all $1 \le K$ $0 \le M$	For all cases: $1 \leq K \leq N \leq 10^9$ $0 \leq M \leq \min(N-1, 200000)$		
	Points	Constraints	
1	3	M = 0	
2	4	K = 1	
3	13	$egin{array}{ll} K=2\ 2\leq N\leq 200000 \end{array}$	
4	24	K = 2	
5	19	$K \leq 100 \ N \leq 200000$	
6	37	No additional constraints	

- M = 0
- No snakes on the board
- Just calcuate the answer by math
- Ans = ceil(n / k) = (n 1) / k + 1

- Time Complexity = O(1)
- Easy 3 points :)

- K = 1
- Each time you can only move 1 square forward
- If there is any snake, which will bring you backward
- It is impossble to go to square N
- Else you can always move to square N by N moves

- Time Complexity = O(1)
- Easy 4 points :)

• Observation : There always exist an optimal solution which won't use the



- The problem become....
- Given a board consist of N+1 squares and M of them are forbidden
- Each time you can move forward 1 K squares
- Find the minimum number of moves required to move from square 0 to square N / determine whether you can do it

- K = 2, 2 <= N <= 200000
- N is small
- Just simulate the process
- Notice that it is optimal to go to the farthest square which is not forbidden

- Let you are now at square i
- If square i+2 is not forbidden, go there
- else if square i+1 is not forbidden, go there
- else return impossible
- Repeat this process until you get to square N

• Time complexity = O(N)

- K = 2, N could be very large, up to 10^9
- We can use the fast forward technique
- For each snake, we calculate the number of moves needed to get over it

- Let you are at square x and the snake entry is at square y
- if the parity of x and y is not same, then you can ignore that snake
- if the parity of x and y is same, then you need to do some simple calculations

- After that, update your position
- After you pass over all M snake, you can calculate the amount of move to go to square N just like what we did in subtask 1
- You can check the impossible case by checking whether there exist 2 consecutive snake

• Time complexity = O(M)

- K <= 100, 2 <= N <= 200000
- K can vary now
- But it is not a big deal

- We can use our algorithm in subtask 3
- check K squares instead of 2 squares

- Let you are now at square i
- If square i+k is not forbidden, go there
- else if square i+k-1 is not forbidden, go there
- else if
- else return impossible
- Repeat this process until you get to square N

• Time complexity = O(N)

- Why time complexity = O(N) instead of O(NK) ?
- although both is fast enough to pass this subtask

- In every 2 moves, you can always move K or more squares forward
- else it is impossible
- O(K) * O(N / K) = O(N)

• You can also use a O(NK) dynamic programming to pass this subtask

- 1 <= K <= N <= 10⁹, 0 <= M <= min(N-1, 200000)
- N could be very large

- Just like what we did before
- We can use the fast forward technique to solve this problem

- Firstly, check the impossible case by checking whether there exist K consecutive snakes
- Then, for the first snake which is in front of you, with position y
- move to the square x' which is one move away from the snake
 - i.e. x' < y, x' + K >= y
- After that, move one more move, find the farthest square which is not forbidden and go there
 - can be done by while loop on the snakes
- If you still can't pass over that snake, i.e. farthest square = y 1, move one more move

- Finally, when there is no snake in front of you, calculate the amount of move to go to square N just like what we did in subtask 1
- Remeber to check the impossible case

• Time complexity = O(M)

• Be careful about the calculation

THX