# J182－Rope 

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## The Problem



ILLEGAL
Cell not covered


OK
3 ropes used


ILLEGAL
Rope covers painted cell


OK
4 ropes used


ILLEGAL Rope overlaps
with itself

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ILLEGAL
Rope overlaps with another rope

## SCORING

| ILLEGAL or $\mathbf{M}>\mathbf{R}+\mathbf{C}+\mathbf{N}$ | $0 \%$ |
| :--- | :--- |
| NICE（i．e．$N+1<\mathrm{M} \leq \mathrm{R}+\mathrm{C}+\mathrm{N})$ | $60 \%$ |
| EXCELLENT（i．e． $\mathrm{M} \leq \mathrm{N}+1$ ） | $100 \%$ |

## SUBTASKS

For all cases：
$1 \leq R, C \leq 300$
$0 \leq N<R \times C$
Points Constraints
1

| 8 | $R=C=2$ |
| :--- | :--- |
|  | $N=0$ |
| 18 | $R=1$ |
| 21 | $N$ |

453 No additional constraints

## Background

## Problem Idea By－percywtc

Testdata By
－percywtc；microtony


Initial version of this problem is a bit harder than the current version This harder version will be discussed later on

## Statistics

0 points $\quad 23+6+1+0=30$
8 points $\quad 7+4+4+0=15$
29 points $\quad 9+9+8+2=28$
47 points $0+1+2+3=6$
100 points $0+0+0+3=3$

First solved by mtyeung1 at 1h 17m 44s

## SUBTASKS

| For all cases： |  |  |
| :---: | :---: | :---: |
| $1 \leq R, C \leq 300$ |  |  |
| $0 \leq N<R \times C$ |  |  |
|  | Points | Constraints |
| 1 | 8 | $R=C=2$ |
|  |  | $N=0$ |
| 2 | 18 | $R=1$ |
| 3 | 21 | $N=0$ |
| 4 | 53 | No additiona |

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## Solution 1 －The First Subtask

8 points for solving a single case $-\mathrm{R}=\mathrm{C}=2, \mathrm{~N}=0$

## INPUT



220

| 1 |  |  | 1 |  |
| :--- | :--- | :--- | :--- | :--- |
| 4 |  |  | 4 |  |
| 1 | 1 | OR | 2 | 1 |
| 1 | 2 |  | 2 | 2 |
| 2 | 2 |  | 1 | 2 |
| 2 | 1 |  | 1 | 1 |


| 1 |  |  |  | 1 |  |
| :--- | :--- | :--- | :--- | :--- | :--- |
| 4 |  |  | 4 |  |  |
| 1 | 1 | OR | 1 | 2 |  |
| 2 | 1 |  | 2 | 2 |  |
| 2 | 2 |  | 2 | 1 |  |
| 1 | 2 |  | 1 | 1 |  |




| 1 |  |  |  |  |  |
| :--- | :--- | :--- | :--- | :--- | :--- |
| 4 |  |  | 4 |  |  |
| 1 | 2 | OR | 2 | 2 |  |
| 1 | 1 |  | 2 | 1 |  |
| 2 | 1 |  | 1 | 1 |  |
| 2 | 2 |  | 1 | 2 |  |

$\begin{array}{llllll}1 & & & & 1 & \\ 4 & & & 4 & \\ 2 & 1 & \text { OR } & 2 & 2 \\ 1 & 1 & & 1 & 2 \\ 1 & 2 & & 1 & 1 \\ 2 & 2 & & 2 & 1\end{array}$

## Solution 1 －The First Subtask

## This solution can only solve Subtask 1 with $\mathbf{M}=\mathbf{1}$ rope used

| Subtask | Score | Max Score |
| :---: | :---: | :---: |
| 1 | 8 | 8 |
| 2 | 0 | 18 |
| 3 | 0 | 21 |
| 4 | 0 | 53 |
| Total | $\mathbf{8}$ | $\mathbf{1 0 0}$ |

## Solution 1 －The First Subtask

## PSEUDOCODE

PrintLine（1）<br>PrintLine（4）<br>PrintLine（1 1）<br>PrintLine（1 2）<br>PrintLine（2 2）<br>PrintLine（2 1）

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## Solution 2 －Dots

Simply，for each unoccupied cell，place a rope

## INPUT



334
13
21
23
32

OUTPUT


5
1
11
1
12

INPUT


337
11
12
13
22

OUTPUT


2
1
21
1
33

## Solution 2 - Dots

In general, this solution needs $\mathbf{M}=\mathbf{R} \times \mathbf{C}-\mathbf{N}$ ropes
Worst case is when $\mathbf{N}=\mathbf{0}$, which requires $\mathbf{M}=\mathbf{R} \times \mathbf{C}$ ropes
For Subtask 1, not worse than NICE (i.e. $\mathrm{M} \leq \mathrm{R}+\mathrm{C}+\mathrm{N}) \quad[4 \leq 2+2+0]$
For Subtask 2, not worse than NICE (i.e. $\mathrm{M} \leq \mathrm{R}+\mathrm{C}+\mathrm{N}) \quad[\mathrm{C} \leq 1+\mathrm{C}+0]$
For Subtask 3 and 4 , may reach $\mathrm{M}>\mathrm{R}+\mathrm{C}+\mathrm{N}$
$[R \times C>R+C+0]$

| Subtask | Score | Max Score |
| :---: | :---: | :---: |
| 1 | 4.8 | 8 |
| 2 | 10.8 | 18 |
| 3 | 0 | 21 |
| 4 | 0 | 53 |
| Total | $\mathbf{1 5 . 6}$ | $\mathbf{1 0 0}$ |

## Solution 2 －Dots

## PSEUDOCODE（For Subtask 2 Only）

```
For i = 1 .. N
    Read(x, y)
    A[y] = True
PrintLine(C - N)
For i = 1 .. C
    If (A[1][i] = False)
    PrintLine(1)
    PrintLine('1 ', i)
```

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## Solution 3 －Snake

For the entire grid empty（ $\mathbf{N}=\mathbf{0}$ in Subtask 3），
We can only use $\mathbf{M}=\mathbf{N + 1}=\mathbf{1}$ rope to fill the whole grid

## INPUT



## OUTPUT



## Solution 3 －Snake

This solution always need $\mathbf{M}=1$ ropes，
And it is only suitable for solving $\mathbf{N}=\mathbf{0}$ cases
For Subtask 1 and 3，it is always an EXCELLENT placement（i．e． $\mathrm{M} \leq \mathrm{N}+1$ ） This solution is not suitable for Subtask 2 and 4

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| Subtask | Score | Max Score |
| :---: | :---: | :---: |
| 1 | 8 | 8 |
| 2 | 0 | 18 |
| 3 | 21 | 21 |
| 4 | 0 | 53 |
| Total | $\mathbf{2 9}$ | $\mathbf{1 0 0}$ |

## Solution 4 －Long Rope in Interval

## For $\mathbf{R}=1$（Subtask 2），

optimal way is for each consecutive unoccupied interval，place a long rope It is easy to see that this is the only optimal way

## INPUT

## OUTPUT

|  |  |  |  |  | $\sum$ | 3 |  | $\square$ |  | $\checkmark$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |

3
13
17
18
110
2
11
12
3
14
15
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## Solution 4 －Long Rope in Interval

This solution needs at most $\mathbf{M}=\mathbf{N}+1$ ropes
as $\mathbf{N}$ obstacles can divide the row into no more than $\mathbf{N}+1$ intervals Worst case is when $\mathbf{N}=\mathbf{0}$

For Subtask 2，must be EXCELLENT（i．e． $\mathrm{M} \leq \mathrm{N}+1$ ）

$$
[N+1 \leq N+1]
$$

This solution is not applicable to Subtask 1， 3 and 4

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| Subtask | Score | Max Score |
| :---: | :---: | :---: |
| 1 | 0 | 8 |
| 2 | 18 | 18 |
| 3 | 0 | 21 |
| 4 | 0 | 53 |
| Total | $\mathbf{1 8}$ | $\mathbf{1 0 0}$ |

## Solutions Summary

| Solutions |  | $\mathbf{1}$－First Sub | 2－Dots | 3－Snakes | 4－Interval |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Subtask | Max Score |  | Score |  |  |
| 1 | 8 | 8 | 4.8 | 8 | 0 |
| 2 | 18 | 0 | 10.8 | 0 | 18 |
| 3 | 21 | 0 | 0 | 21 | 0 |
| 4 | 53 | 0 | 0 | 0 | 0 |
| Total | $\mathbf{1 0 0}$ | $\mathbf{8}$ | $\mathbf{1 5 . 6}$ | $\mathbf{2 9}$ | $\mathbf{1 8}$ |


| Score |
| :---: |
| 8 |
| 18 |
| 21 |
| 0 |
| 47 |

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## Solution 4 －Long Rope in Interval

## PSEUDOCODE（For Subtask 2 Only）

Count $=0$
$A[0]=A[C+1]=$ True
For $i=0 \ldots C$
If（A［i］AND A［i＋1］） Count＋＋
PrintLine（N＋ 1 －Count）

```
For i = 0 .. C
    If (A[i] AND NOT(A[i+1]))
        j = i + 1
        While (j<C+1 AND A[j]=False)
        j++
    PrintLine(j - i)
    For k = i .. j
        PrintLine('1 ', k)
    i = k + 1
```

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## Solution 5 －Copies of Solution 4

We can consider the each row separately
For each row，simply apply Solution 4 once

## INPUT



## OUTPUT

|  |  | \％ |  |  |  | x | $\times$ |  | x |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| ＊ |  | － | $\times$ | $\times$ |  |  |  |  |  |
|  |  | － | $x$ | $\times$ | $\times$ | ， |  | ＊ |  |

## Solution 5 －Copies of Solution 4

This solution needs at most $\mathbf{M}=\mathbf{R +} \mathbf{N}$ ropes
$\mathbf{n}$ obstacles can divide a row into no more than $\mathbf{n + 1}$ intervals
So，All $\mathbf{R}$ rows have in total no more than $\mathbf{N}+\mathbf{R}$ intervals
For Subtask 2，must be EXCELLENT（i．e． $\mathrm{M} \leq \mathrm{N}+1$ ）
$[\mathrm{N}+1 \leq \mathrm{N}+1]$
For Subtask 1， 3 and 4，ast least NICE（i．e．$M \leq R+C+N) \quad[N+R \leq R+C+N]$

| Subtask | Score | Max Score |
| :---: | :---: | :---: |
| 1 | 4.8 | 8 |
| 2 | 18 | 18 |
| 3 | 12.6 | 21 |
| 4 | 31.8 | 53 |
| Total | $\mathbf{6 7 . 2}$ | $\mathbf{1 0 0}$ |

## Solutions Summary



## Solution 6 －Mix of Solution 3 \＆ 4

We can consider the snake in Solution 3 as a long line Using Solution 4 －Long Rope in Interval on this long line

INPUT


PROCESS


OUTPUT


## Solution 6 - Mix of Subtask 3 \& 4

This solution considers a single long line with $\mathbf{N}$ obstacles
As mentioned in Solution 4, this requires at most $\mathbf{M}=\mathbf{N}+1$ ropes
For all Subtasks, it is always EXCELLENT (i.e. $\mathrm{M} \leq \mathrm{N}+1$ )
$[N+1 \leq N+1]$

| Subtask | Score | Max Score |
| :---: | :---: | :---: |
| 1 | 8 | 8 |
| 2 | 18 | 18 |
| 3 | 21 | 21 |
| 4 | 53 | 53 |
| Total | $\mathbf{1 0 0}$ | $\mathbf{1 0 0}$ |

## Summary

This task does not require you to solve every case with optimal arrangement， But just＂good enough＂way

It is always a good idea to think of some small or special cases at first It may lead you to develop a better solution

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## Other Solution－Depth First Search（DFS）

Choose an unoccupied cell， Repeatedly walking to any unoccupied cell next to it until cannot do so Repeat the whole process with another rope if there are still unoccupied cells
The score depends on your implementation：

| Subtask | Score | Max Score |  | Subtask | Score | Max Score |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 8 | 8 |  | 1 | 8 | 8 |
| 2 | 18 | 18 |  | 2 | 18 | 18 |
| 3 | 21 | 21 |  | 3 | 21 | 21 |
| 4 | 31.8 | 53 |  | 4 | 0 | 53 |
| Total | $\mathbf{7 7 . 8}$ | $\mathbf{1 0 0}$ |  | Total | $\mathbf{4 7}$ | $\mathbf{1 0 0}$ |

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## Harder Version of the Problem

## CONSTRAINTS

$\mathrm{R}, \mathrm{C} \geq 2$
$\mathbf{R} \times \mathbf{C}$ is even
$\mathrm{N} \geq 1$
Only $\mathbf{N}$ ropes can be used

Can you solve it？：）

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