

HKOI 2016/17
JQ3 – Fibonacci Word

Alex Tung

21/1/2017

Problem description

- $F(1) = \text{"0"}, F(2) = \text{"01"}$
- $F(n+2) = F(n+1) + F(n)$
- $F = \text{"010010100100101001010..."}$
- Find number of occurrences of W_i in $F[L..R]$
- $W_i = \text{"00"} / \text{"01"} / \text{"10"} / \text{"11"}$

SUBTASKS

For all cases: $1 \leq Q \leq 10000$, $1 \leq L_i \leq R_i \leq 10^{18}$

	Points	Constraints
<i>1</i>	3	$W_i = 11$
<i>2</i>	24	$R_i \leq 2000$
<i>3</i>	17	$R_i \leq 1000000$
<i>4</i>	19	$R_i - L_i \leq 100$
<i>5</i>	37	No additional constraints

Statistics

Attempts	Max	Mean	Std Dev	Subtasks				
65	100	14.015	21.09	3: 53	24: 21	17: 8	19: 2	37: 2

Subtask 1

- $W = \text{"11"}$
- Prove / Observe that $\text{answer} = 0 :P$

Subtasks 2 – 3 (R small)

- Generate a prefix of F
- After 28 iterations,
F.length() = 1346269

```
//C++ implementation

string F, temp, temp2;

F = "01";
temp = "0";

for(int i = 1; i <= 28; i++){
    temp2 = F + temp;
    temp = F;
    F = temp2;
}
```

$O(RQ)$ solution

- Step 1: generate string
- Step 2: for each query, count **directly**

$O(Q)$ solution

- Step 1: generate string
- Step 2: **precompute** the number of occurrences of “00” / “01” / “10” / “11” in $F[1..K]$
- Step 3: for each query (L, R, W) ,
output $\text{cnt}[W][R] - \text{cnt}[W][L]$

(This technique is called ‘partial sum’.)

Example

- $F = \text{"01\underline{00}101\underline{00}1\underline{00}101\underline{00}1010\dots\text{"}$
- $\text{query}(4, 12, \text{"00"})$

i	1	2	3	4	5	6	7	8	9	10	11	12
$\text{cnt}[\text{"00"}][i]$	0	0	0	<u>1</u>	1	1	1	1	2	2	2	<u>3</u>

- $\text{Answer} = \text{cnt}[\text{"00"}][12] - \text{cnt}[\text{"00"}][4]$
 $= 3 - 1$
 $= 2$

Subtask 4 ((R - L) small)

- Idea: want to generate $F[L\dots R]$, by finding out what **each character is**

`len[n] := length of F(n)`



`pos -= len[i+1];`

- Time complexity: $O(Q \log R)$

Subtask 5

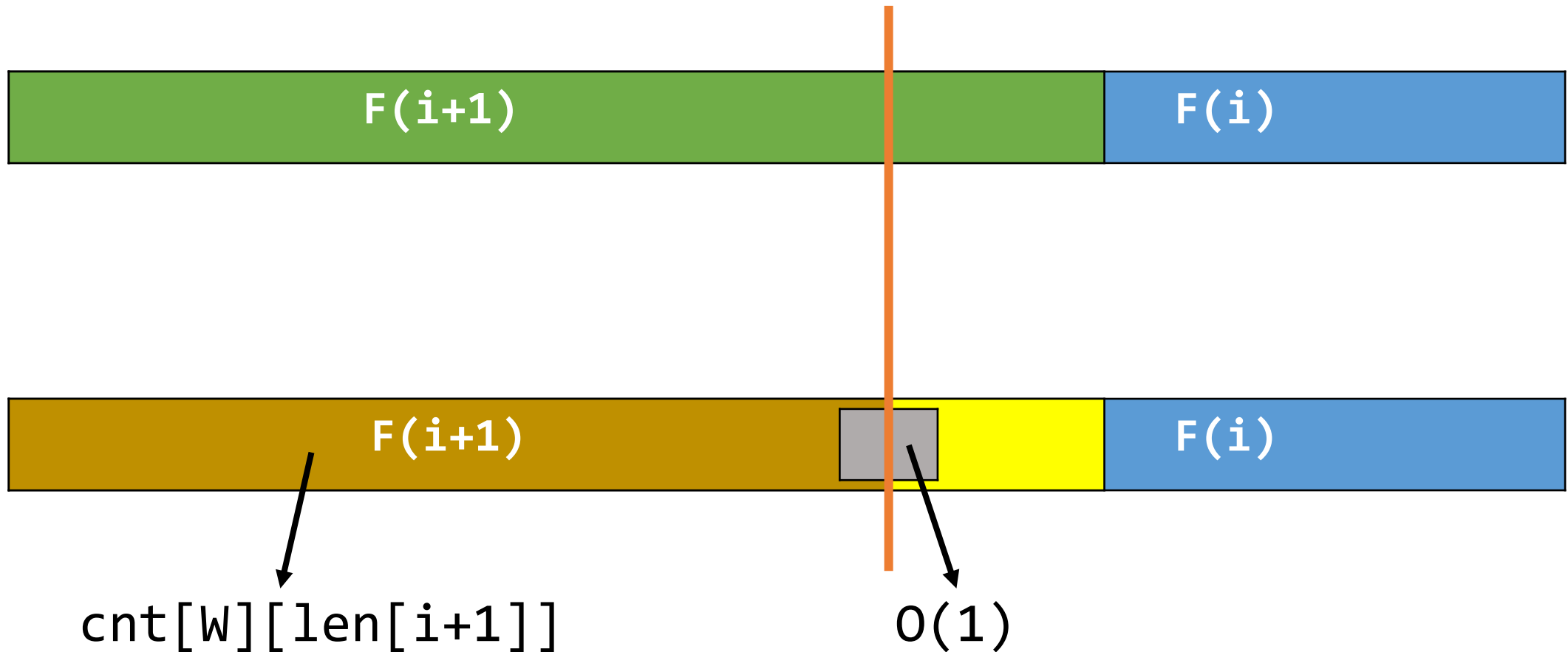
- Merge ideas from subtask 3 solution and subtask 4 solution
- Idea 1: “`answer = cnt[W][R] - cnt[W][L]`”
- Idea 2: “`pos -= len[i+1];`”

The algorithm

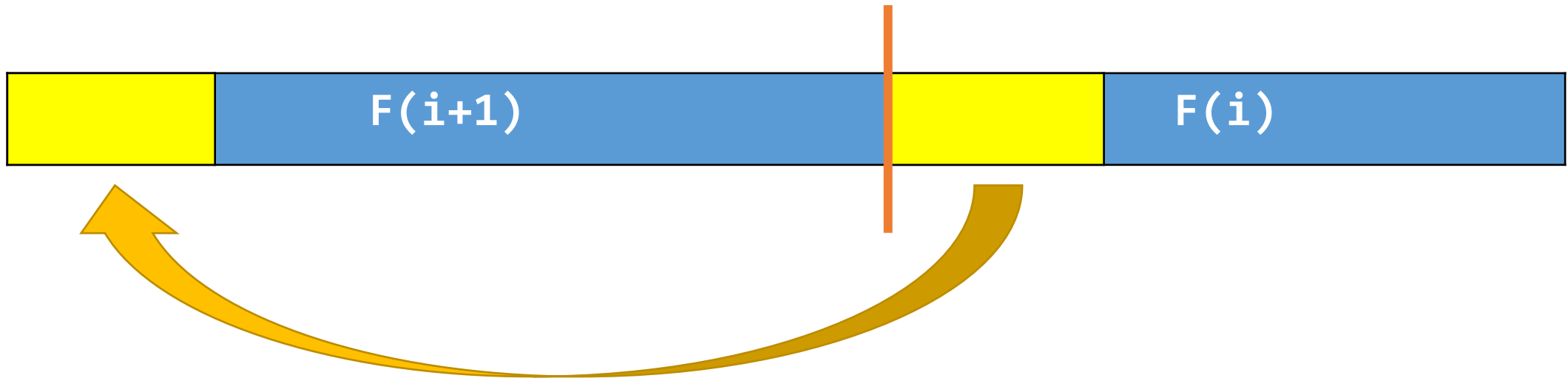
- Step 1: Find the values of `cnt[W][len[i]]` and the first and last characters of `F(i)`
- First character: always '0'
- Last character: alternating between '0' and '1'

- Step 2: Calculate $\text{query}(1, R, W)$ by **reducing it into smaller cases** (similar for $\text{query}(1, L, W)$)

Calculating query(1, R, W)



Calculating query(1, R, W)



- Time complexity: $O(Q \log R)$
- To solve for $R \leq 10^{18}$, need to find `cnt[W][len[1]]...cnt[W][len[87]]`

The End

- Any questions?