

MISC CS TOPICS (IV)

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Generations of Programming Language

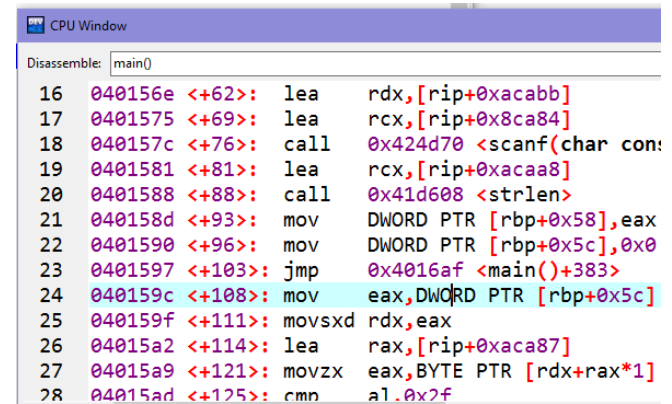
	1 st Generation Machine Code	2 nd Generation Assembly	3 rd Generation High-level	4 th Generation Declarative
How to run	Can be directly run by CPU	Requires assembler	Requires compiler / interpreter	Requires interpreter
Portability	Depends on CPU / microarchitecture	Depends on architecture	High; Depends on compiler	High; Depends on interpreter
Languages			Pascal C/C++	SQL, Haskell
Example	0x24a60004 0x0001d821 0xac240204	lea eax,[rbp-0x60] sub eax,0x30 jmp 0x4016ab	int x = 0; while (n) { x += n % 10; n /= 10; }	SELECT count(*) FROM 'users' WHERE 'name' LIKE 'CHAN%'

Program and ISA

- A program is an ordered list of **instructions**
- Instructions are commands to the CPU
 - Perform calculations
 - Read data from / write data to memory
 - Control system components
- A **Instruction Set Architecture** defines
 - What instructions are available
 - What data types are available
 - How memory are addressed
 - I/O, Exception handling, Interrupt

Assembly Language

- Low level language that is tightly coupled to a specific ISA
- Requires **assembler** to
 - Detect issues
 - Assign / convert addresses to symbols / labels
 - Convert pseudo-instructions into actual instructions
 - Convert instructions into binary
- Examples of ISAs
 - **x86**, **ARM**, **AVR**, **PowerPC**, **8051**, **MIPS**
- A **disassembler** converts machine code back to assembly code

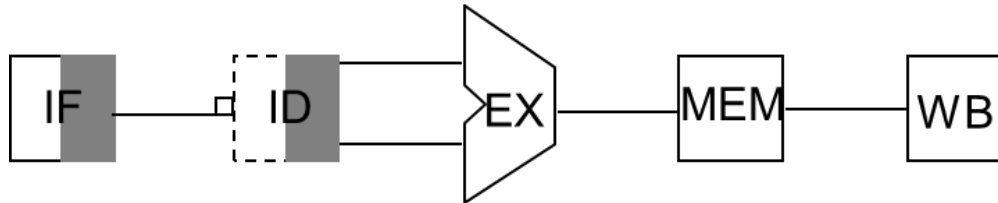


The screenshot shows a window titled "CPU Window" with a "Disassemble:" field containing "main0". Below this, a list of assembly instructions is displayed with their addresses and offsets. The instruction at address 040159c is highlighted in blue.

```
16 040156e <+62>: lea    rdx,[rip+0xacabb]
17 0401575 <+69>: lea    rcx,[rip+0x8ca84]
18 040157c <+76>: call  0x424d70 <scanf(char con:
19 0401581 <+81>: lea    rcx,[rip+0xaca8]
20 0401588 <+88>: call  0x41d608 <strlen>
21 040158d <+93>: mov   DWORD PTR [rbp+0x58],eax
22 0401590 <+96>: mov   DWORD PTR [rbp+0x5c],0x0
23 0401597 <+103>: jmp   0x4016af <main()+383>
24 040159c <+108>: mov   eax,DWORD PTR [rbp+0x5c]
25 040159f <+111>: movsxd rdx,eax
26 04015a2 <+114>: lea   rax,[rip+0xaca87]
27 04015a9 <+121>: movzx eax,BYTE PTR [rdx+rax*1]
28 04015ad <+125>: cmp   al,0x2f
```

MIPS

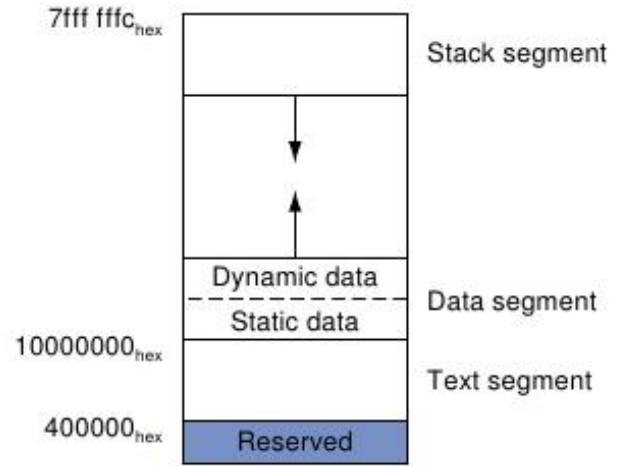
- Microprocessor without Interlocked Pipeline Stages
- MIPS was introduced in 1981
- We'll be using CPU "R2000" (introduced in 1985)
- 32 32-bit General Purpose **registers** (寄存器)
- Instructions are 32-bit too
 - In contrast, x86 instructions have variable length



helloworld.s

Assemblers, Linkers, and the SPIM Simulator
James R. Larus Microsoft Research

```
1      .data
2 msg:  .asciiz "Hello World\n"
3
4      .text
5 main: li $v0, 4      # syscall 4 (print_str)
6      la $a0, msg     # argument: string
7      syscall        # print the string
8
9      jr $ra         # return
```



- Data segment: store constant data (e.g. strings)
reserve space for variables / arrays
- Text segment: actual program
- The program will start from the `main:` label

SPIM (MIPS Simulator)

Program counter
(which instruction runs next)

EVERY TIME You want to run a program, press
Reinitialize and Load File, then Run/Continue

FP Regs	Int Regs [16]	Data	Text
Int Regs [16]		Text	
PC	= 400028	User Text Segment [00400000]..[00440000]	
EPC	= 0	[00400000]	8fa40000 lw \$4, 0(\$29) ; 183: lw \$a0 0(\$sp) # argc
Cause	= 0	[00400004]	27a50004 addiu \$5, \$29, 4 ; 184: addiu \$a1 \$sp 4 # argv
BadVAddr	= 0	[00400008]	24a60004 addiu \$6, \$5, 4 ; 185: addiu \$a2 \$a1 4 # envp
Status	= 3000ff10	[0040000c]	00041080 sll \$2, \$4, 2 ; 186: sll \$v0 \$a0 2
HI	= 0	[00400010]	00c23021 addu \$6, \$6, \$2 ; 187: addu \$a2 \$a2 \$v0
LO	= 0	[00400014]	0c100009 jal 0x00400024 [main] ; 188: jal main
R0 [r0]	= 0	[00400018]	00000000 nop ; 189: nop
R1 [at]	= 0	[0040001c]	3402000a ori \$2, \$0, 10 ; 191: li \$v0 10
R2 [v0]	= 4	[00400020]	0000000c syscall ; 192: syscall # syscall 10 (exit)
R3 [v1]	= 0	[00400024]	34020004 ori \$2, \$0, 4 ; 5: li \$v0, 4 # syscall 4 (print_str)
R4 [a0]	= 1	[00400028]	3c041001 lui \$4, 4097 [msg] ; 6: la \$a0, msg # argument: string
R5 [a1]	= 7ffff354	[0040002c]	0000000c syscall ; 7: syscall # print the string
R6 [a2]	= 7ffff35c	[00400030]	3402000a ori \$2, \$0, 10 ; 9: li \$v0, 10 # syscall 10 (exit)
R7 [a3]	= 0	[00400034]	0000000c syscall ; 10: syscall # exit the program

Registers
("variables" visible by CPU)

Memory Machine
Address Code

Processed
Code

Your Assembly Code

Data segment

- MIPS Assembler provides various directives to help you describe data easily

str1: .asciiz "abcdef\n"

- Creates a null-terminated string with data "abcdef\n"
- The memory location of the string is labelled with str1

```
char str1[] = "abcdef\n";
```

str2: .space 30

```
char str2[30];
```

primes: .word 2, 3, 5, 7, 11, 13

- Creates an integer array of length 6, initialized with 2, 3, 5, 7, 11, 13
- The memory location of the array is labelled with primes

```
int primes[] = {2, 3, 5, 7, 11, 13};
```

array1: .word 0:10

- Creates an array of length 10
- Each element is 4 bytes and is initialized to 0
- The memory location of the array is labelled with array1

```
int array1[10] = {0};
```


Registers

- Registers are like “variables”
- Registers in MIPS are like 32-bit integers
- Some of them have special meaning
 - Arguments to function calls
 - Return value of function
 - Stack counter
- You can use \$t0 - \$t9, \$s0 - \$s7 freely
- \$zero is a special register that is always 0
- Excess data has to be stored in memory

Assignment and addition

li \$rd, imm (pseudo-instruction)

- Load Immediate

- Example: **li \$t0, 5**

Pascal

\$t0 := 5

C / C++

\$t0 = 5

move \$rd, \$rs

- Example: **move \$t0, \$t1**

\$t0 := \$t1

\$t0 = \$t1

add \$rd, \$rs, \$rt

- Adding \$rs and \$rt and store result into \$rd

- Example: **add \$t0, \$t1, \$t2**

\$t0 := \$t1 + \$t2

\$t0 = \$t1 + \$t2

addi \$rd, \$rs, imm

- Adding \$rs and 16-bit signed imm and store result into \$rd

- Example: **add \$t0, \$t1, 1**

\$t0 := \$t1 + 1

\$t0 = \$t1 + 1

Syscalls

- Syscalls allows you to perform input and output
- Controlled by value stored in register \$v0

Service	Code (\$v0)	Arguments	Result
print_int	1	\$a0 = integer	
print_float	2	\$f12 = float	
print_double	3	\$f12 = double	
print_string	4	\$a0 = string	
read_int	5		integer (in \$v0)
read_float	6		float (in \$f0)
read_double	7		double (in \$f0)
read_string	8		\$a0 = buffer, \$a1 = length
sbrk	9	\$a0 = amount	address (in \$v0)
exit	10		
print_char	11	\$a0 = char	
read_char	12		char (in \$a0)

aplusb.s

Console

```
5
9
14
|
```

```
1  .data
2
3  .text
4  main: li $v0, 5      # syscall 5 (read_int)
5        syscall      # read integer
6        # at this point the input is stored in $v0
7        move $t0, $v0 # $t0 = $v0
8
9        li $v0, 5     # syscall 5 (read_int)
10       syscall      # read integer
11
12       add $a0, $t0, $v0 # $a0 = $t0 + $v0
13       li $v0, 1      # syscall 1 (print_int)
14       syscall      # $a0 is printed to screen
15
16       li $a0, 10     # ascii code of end line character
17       li $v0, 11     # syscall 11 (print_char)
18       syscall
19
20       jr $ra        # return
```

Pascal

```
readln($v0);
$t0 := $v0;
```

```
readln($v0);
$a0 := $t0 + $v0;
```

```
write($a0);
```

```
write(chr(10));
```

C / C++

```
scanf("%d", $v0);
$t0 = $v0;
```

```
scanf("%d", $v0);
$a0 = $t0 + $v0;
```

```
printf("%d", $a0);
```

```
printf("%c", 10);
```

```
return 0;
```

Branch instructions

- Instructions normally execute from top to bottom
- To write “if” statements, we need branch instructions

j label

- Unconditional jump
- Example: `j loop`

Pascal

`goto loop;`

C / C++

`goto loop;`

b?? \$rs, label

- Conditional jump (if true)
- Example: `beqz $t0, loop`

`if ($t0 = 0) then`

`goto loop;`

`if ($t0 == 0)`

`goto loop;`

b?? \$rs, \$rt, label (pseudo-instruction)

- Conditional jump (if true)
- Example: `bgt $t0, $t1, loop`

`if ($t0 > $t1) then`

`goto loop;`

`if ($t0 > $t1)`

`goto loop;`

phonenumber.s

```
1      .data
2 fixed: .asciiz "Fixed\n"
3 mobile: .asciiz "Mobile\n"
4
5      .text
6 main: li $v0, 5      # syscall 5 (read_int)
7      syscall        # read integer
8      # at this point the input is stored in $v0
9
10     la $a0, fixed   # assume its Fixed
11     li $t0, 40000000 # $t0 = 40000000
12     blt $v0, $t0, output # if ($v0 < $t0) goto output;
13     la $a0, mobile # else: load mobile instead
14 output:
15     li $v0, 4      # syscall 4 (print_string)
16     syscall        # string pointed by $a0 is
17     # printed to screen
18     jr $ra        # return
```

la \$rd, label

- Load Address
- Example: `la $t0, a`
- Pascal: `var $t0: Pointer;`
`$t0 := @a;`
- C/C++: `void* $t0 = &a;`
- Note that “false” part is above the “true” part

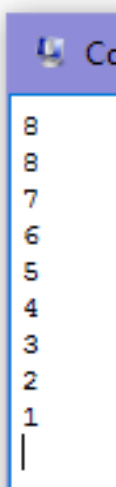
print_string syscall (\$v0 = 4)

- C/C++: `printf("%s", a0);`

Loops

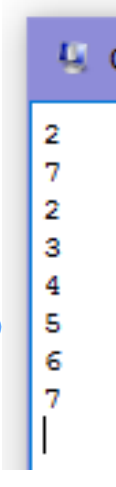
- do-while loop (dowhile.s)

```
1      .data
2
3      .text
4 main:
5      li $v0, 5
6      syscall
7      move $t0, $v0      # $t0 = $v0
8
9 loop:      # do {
10     li $v0, 1          # syscall 1 (print_int)
11     move $a0, $t0      # $a0 = $t0
12     syscall
13
14     li $v0, 11         # syscall 11 (print_char)
15     li $a0, 10        # end of line character
16     syscall
17
18     addi $t0, $t0, -1  # $t0 = $t0 - 1
19     bgt $t0, $zero, loop # } while ($t0 > 0)
20
21     jr $ra            # return
```

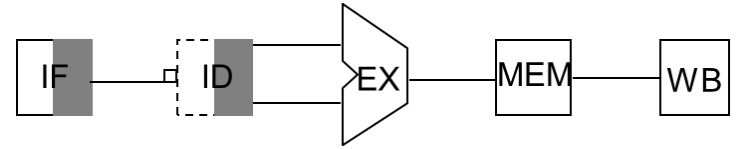


- for loop (forloop.s)

```
1      .data
2
3      .text
4 main:
5      li $v0, 5
6      syscall
7      move $t1, $v0      # $t1 is left
8      li $v0, 5
9      syscall
10     move $t2, $v0      # $t2 is right
11
12     move $t0, $t1      # int $t0 = $t1
13 loop:
14     bgt $t0, $t2, abc  # ; $t0 <= $t2;
15     li $v0, 1          # syscall 1 (print_int)
16     move $a0, $t0      # $a0 = $t0
17     syscall
18
19     li $v0, 11         # syscall 11 (print_char)
20     li $a0, 10        # end of line character
21     syscall
22
23     addi $t0, $t0, 1   # $t0 = $t0 + 1
24     j loop            # unconditional jump
25
26 abc: jr $ra          # return
```



Accessing Memory



- Unlike x86 and many other architectures MIPS is a load-store architecture
- Only the following instructions can access memory:
 - Load: **lw lh lb** Store: **sw sh sb**
 - **w h b** stands for 4 bytes, 2 bytes and 1 byte respectively
- Bare machine provides only one addressing mode: $c(rx)$, which uses the sum of the immediate c and register rx as the address
- Addressing formats allowed by SPIM

Format	Address computation
(register)	contents of register
imm	immediate
imm (register)	immediate + contents of register
label	address of label
label \pm imm	address of label + or - immediate
label \pm imm (register)	address of label + or - (immediate + contents of register)

Strings

- syscall 8 (read_string) is similar to C function `fgets($a0, $a1, stdin)`
 - \$a0 is the memory address of the buffer (where the string can be stored to)
 - \$a1 is the length of the buffer

lb \$rd, offset(\$rs)

- Load Byte (8-bit)
- Example: `lb $t0, 5($t1)`
- Pascal (t0 is char, t1 is ^char): `$t0 := ($t1 + 5)^;`
- C/C++ (t0 is char, t1 is char*): `$t0 = $t1[5];`

```
1      .data
2  str1: .space 30
3
4      .text
5  main:
6      li $v0, 8
7      la $a0, str1      # address of buffer
8      li $a1, 30        # size of buffer
9      syscall
10     # at this point the line is stored in address $a0
11
12     la $s0, str1
13
14     li $v0, 1          # syscall 1 (print_int)
15     lb $a0, 0($s0)    # get ascii code of $s0[0]
16     syscall
17
18     li $v0, 11         # syscall 11 (print_char)
19     li $a0, 10        # end of line character
20     syscall
21
22     li $v0, 1          # syscall 1 (print_int)
23     lb $a0, 2($s0)    # get ascii code of $s0[2]
24     syscall
25
26     li $v0, 11         # syscall 11 (print_char)
27     li $a0, 10        # end of line character
28     syscall
29
30     jr $ra
```

readstring.s

Co
ABCD
65
67
|

Arrays

- Integers are 4-bytes wide
- To access an integer array, multiply index by 4 add then base address

\$t2=

0x1000	0x1004	0x1008	0x100C	0x1010	0x1014
2	3	5	7	11	13

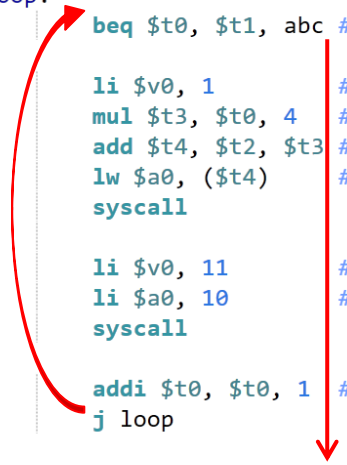
lw \$rd, offset(\$rs)

- Load Word (32-bit)
- Example: `lw $t0, 8($t1)`
- Pascal (t0 is longint, t1 is ^longint): `$t0 := ($t1 + 2)^;`
- C/C++ (t0 is int, t1 is int*): `$t0 = $t1[2];`

sb \$rd, offset(\$rs) Store Byte

```
1      .data
2 primes: .word 2, 3, 5, 7, 11, 13
3
4      .text
5 main:
6      li $t0, 0      # int i = 0;
7      li $t1, 6
8      la $t2, primes
9 loop:
10     beq $t0, $t1, abc # ; $t0 != $t1;
11
12     li $v0, 1      # syscall 1 (print_int)
13     mul $t3, $t0, 4 # $t3 = $t0 + 4 (number of bytes)
14     add $t4, $t2, $t3 # $t4 = $t2 + $t3 (memory address)
15     lw $a0, ($t4)  # get $t4[0]
16     syscall
17
18     li $v0, 11     # syscall 11 (print_char)
19     li $a0, 10     # end of line character
20     syscall
21
22     addi $t0, $t0, 1 # $t0 = $t0 + 1
23     j loop
24
25 abc: jr $ra
```

primes.s



2
3
5
7
11
13
|

sw \$rd, offset(\$rs) Store Word

Pseudo instructions

- Instruction is only 32 bits. It is impossible load large values. Separate into two instructions (largenumber.s)

```
li $t0, 20000000      [00400024] 3c010131 lui $1, 305           ; 5: li $t0, 20000000
li $t1, 30000000      [00400028] 34282d00 ori $8, $1, 11520    ; 6: li $t1, 30000000
                        [0040002c] 3c0101c9 lui $1, 457           ; 6: li $t1, 30000000
li $v0, 1              [00400030] 3429c380 ori $9, $1, -15488 ; 8: li $v0, 1
add $a0, $t0, $t1     [00400034] 34020001 ori $2, $0, 1       ; 8: li $v0, 1
syscall                [00400038] 01092020 add $4, $8, $9       ; 9: add $a0, $t0, $t1
                        [0040003c] 0000000c syscall          ; 10: syscall
```

- Array access (primes2.s)

```
12      ..... mul $t3, $t0, 4 # $t3 = $t0 + 4 (number of bytes)
13      ..... lw $a0, primes + 0($t3) # get primes + $t3 bytes
[00400034] 34010004 ori $1, $0, 4           ; 12: mul $t3, $t0, 4 # $t3 = $t0 + 4 (number of bytes)
[00400038] 71015802 mul $11, $8, $1
[0040003c] 3c011001 lui $1, 4097           ; 13: lw $a0, primes + 0($t3) # get primes + $t3 bytes
[00400040] 002b0821 addu $1, $1, $11
[00400044] 8c240000 lw $4, 0($1)
```

Next Week

- Fast Forward 35 years
- Learn the most advanced technology today
- Write much faster programs
- Mini-comp again!