

# HKOI 2015/16 TFT Q1 Solution

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# Comments

- Easy-Medium
- Greedy algorithm; less standard
- More full scorers (5) than expected!
- Lack of partial solution for  $K \leq 10$  is puzzling

# Observation 1

- For each type  $X$  of apple, we only need to consider
  - the first tall tree with type  $X$  apples and
  - the first short tree with type  $X$  apples
- Plus, we take exactly one apple of each type (thus, in all cases where winning is possible, number of trees climbed =  $K$ )

# Observation 2

- $N$  is irrelevant
- For each tree, we only need to know the number of the first checkpoint after it

# Notations

For  $1 \leq i \leq K$ ,

- Let  $T[i]$  be the first checkpoint after the first tall tree with type  $i$  apples (-1 if not exist)
- Let  $S[i]$  be the first checkpoint after the first short tree with type  $j$  apples (-1 if not exist)

# Example

- $(N = 5, K = 2)$
- T 1, T 2, T 1, S 1, T 1
- 2 5
  
- $T[1] = 1; S[1] = 2$
- $T[2] = 1; S[2] = -1$

- Note that the notation used in this presentation is only convenient for demonstrating the idea; it complicates the process of retrieving the answer

# Solution 1

- Subtask 1 (all trees are short)
- $T[X] = -1$  for all  $X$

```
for i from 1 to K
```

```
  for j from 1 to K
```

```
    if (type j apple has not been taken yet) and ( $S[X] \leq i$ )
```

```
      take type j apple for the i-th checkpoint
```

```
      break
```



# Solution 1

- Time complexity:  $O(K^2)$
- Expected score: 8

# Solution 2

- Subtask 2 ( $N = K+1$ )
- We need at least  $K$  apples; that means we can only give up climbing one tree!
- Exhaustion on which tree to not climb
- For each case, simulate the race

# Solution 2

- Time complexity:  $O(K^2)$
- Expected score: 15
- Expected score (combined with solution 1): 23

# Solution 3

- Solves for  $K \leq 10$
- Thanks to observation 1, we can exhaust, for each type of apple, whether to climb **the tall tree** or **the short tree**
- Then, check whether RB can pass all checkpoints in  $O(K)$  /  $O(K^2)$

# Solution 3

- Time complexity:  $O(2^K * K)$  /  $O(2^K * K^2)$
- Expected score: 27
- Expected score (combined with solutions 1+2): 50

# Solution 4 (greedy)

- We consider the checkpoints one by one
- For each checkpoint, we want to choose the “best” type of apple among those available for choice

# Solution 4

- Climbing a short tree is always more desirable than climbing a tall tree
- If there are multiple available short trees, any one will do
- If no short trees are available, climb the tall tree  $T[X]$  such that  $S[X]$  is maximal
- Intuitively, this is because it takes the longest time to 'wait' for  $S[X]$  to be available

# Solution 4

```
for i from 1 to K
  //try short tree
  for j from 1 to K
    if (!used[j]) and (S[j] <= i)
      climb(S[j]) //perform necessary updates
```



# Solution 4

```
//try tall tree; only if no short trees are available
maxi = -1
choose = -1
for j from 1 to K
    if(!used[j]) and (T[j] <= i) and S[j] > maxi
        choose = j
        maxi = S[j]
climb(T[choose])
//if still no trees, 'Lose'
```

# Solution 4

- Time complexity:  $O(N + K^2)$
- Expected score: 80
  
- P.S. this task was originally proposed for HKOI 2015/16, without the last subtask (i.e. solution 4 would suffice to get 100 points)

# Solution 5

- We speed up solution 4

Moving from one checkpoint to the next,

- we add the newly available tall trees,  $T[i]$ , into a heap (e.g. C++ `priority_queue`), using  $S[i]$  for ordering
- we add the newly available short trees,  $S[i]$ , into a queue (many other data structures will also do)

# Solution 5

- Then, do the same as in solution 4
- Short tree is better than tall tree
- Tall tree with larger  $S[X]$  is better than tall tree with smaller  $S[X]$

# Solution 5

- Time complexity:  $O(N + K \log K)$
- Expected score: 100 :)