

# Competition

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# Roadmap

- ▶ Problem
- ▶ Statistics
- ▶ 50% solution
- ▶ Greedy solution

# Problem

- ▶ 田忌賽馬
- ▶ Two teams of students take part in a one-on-one competition
- ▶ Find the maximum number of rounds which Team A can win

# Statistics

- ▶ 15 full marks
- ▶ Highest mean
- ▶ Highest standard deviation

# 50% solution

- ▶ There are  $N!$  ways of assigning students from Team A
- ▶ Enumerate them one by one and count the number of winning rounds
- ▶ When  $N \leq 10$ , it takes at most  $10! = 3628800$  iterations
- ▶ Implementation: recursion
- ▶ Enough to score 50 marks

# Greedy solution

## Idea

Keep assigning the strongest student from Team A to beat the student from Team B who is just weaker than him/her

## Algorithm

1. Find the strongest student from Team A, call him/her  $A_i$
2. Find the student from Team B who is just weaker than  $A_i$ , call him/her  $B_j$
3. If  $B_j$  does not exist, terminate
4. Increase counter, remove  $A_i$  and  $B_j$  from consideration, iterate

# Proof of correctness

Let  $A_i$  be the strongest student from Team A

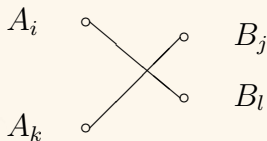
Let  $B_j$  be the student from Team B who is just weaker than  $A_i$

Assume in an optimal assignment  $S$ ,  $A_i$  competes with  $B_l$  and

$A_k$  competes with  $B_j$

Further assume  $B_l$  and  $B_j$  are different students

Optimal assignment  $S$



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New assignment  $S'$

$A_i$          $B_j$

$A_k$          $B_l$

What happens if  $A_i$  and  $A_k$  exchange their competitors?



## First case: $A_i > B_l$

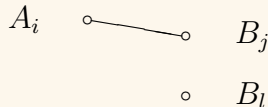
- ▶  $A_i > B_l$  means  $A_i$  is stronger than  $B_l$
- ▶ By swapping competitors,  $A_i$  can still win a round in  $S'$
- ▶  $A_k$  now competes with  $B_l \leq B_j$ , and the result will not be worse

Optimal assignment  $S$



$\Rightarrow$

New assignment  $S'$



## Second case: $A_i \leq B_l$

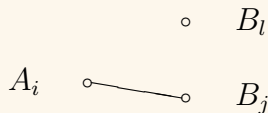
- ▶  $A_i$  and  $A_k$  together win at most one round in  $S$
- ▶  $A_i$  and  $A_k$  together win at least one round in  $S'$

Optimal assignment  $S$



$\Rightarrow$

New assignment  $S'$



# Conclusion

- ▶ There always exists an optimal assignment in which the strongest student from Team A competes with the one from Team B who is just weaker
- ▶ Clearly an  $O(N^2)$  implementation is possible
- ▶  $O(N \log N)$  implementations also exist
- ▶ Keep it simple, stupid (KISS)