Hong Kong Olympiad in Informatics 2016/17
Heat Event (Junior Group)
Official Solution

## Statistics ( $\mathbf{N}=\mathbf{2 5 6}$ )

Full mark $=45$. Maximum $=41$. Median $=16.75$. Advance to Final $=20$ marks or above.

## Section A

| Q | A | Explanation |
| :---: | :---: | :---: |
| 1 | F | Random-Access Memory(RAM) is a hardware which cannot be downloaded from the Internet. |
| 2 | F | Both USB 2.0 and USB 3.0 data signals are digital. One of the reason why USB 3.0 has higher data transfer rate is because it supports full duplex data transfer. |
| 3 | T | A char requires 1 byte ( 8 bits) to represent $0 \sim 255$ (or $-128 \sim 127$ depend on compiler). |
| 4 | T | No matter when x is true or false, ( x and (not x$)$ ) is always false. |
| 5 | F | Here is the proof. |
| 6 | D | Ransomware is a type of malicious software which encrypts the victim's files, making them inaccessible and demands a ransom payment to decrypt them. |
| 7 | A | The nature of nested function calling and returning is the same as that of a stack. |
| 8 | B | $\begin{aligned} & \mathrm{A}=[0,6] \\ & \mathrm{B}=[0,4] \\ & \mathrm{B}-\mathrm{A}=[0,4]-[0,6]=[-6,4]=11 \text { possible outputs. } \end{aligned}$ |
| 9 | B | Assume B is the correct answer. <br> Let the correct answer be T and the wrong answer be F. Since there is only one correct answer: $\mathrm{A}=\mathrm{F}, \mathrm{~B}=\mathrm{T}, \mathrm{C}=\mathrm{F}, \mathrm{D}=\mathrm{F}$ <br> When A is F , $(\mathrm{B}!=\mathrm{T}$ and $\mathrm{C}!=\mathrm{T})$ need to be false <br> When B is T , $(\mathrm{C}!=\mathrm{T}$ and $\mathrm{D}!=\mathrm{T})$ need to be true <br> When C is $\mathrm{F},(\mathrm{C}=\mathrm{T})$ need to be false <br> When D is F , ( $\mathrm{A}!=\mathrm{T}$ and $\mathrm{B}!=\mathrm{T}$ ) need to be false. <br> The situation in above satisfy all the restrictions. So B is the answer, other cases will result in contradiction. |
| 10 | C | AlphaGo is a Ai computer program developed by Google that plays the board game Go. <br> Deep blue is designed for playing chess. Pokémon Go is a mobile phone game. DuckDuckgo is a search engine. |

11 D Program A is insertion sort as in each for loop, a[0..i-1] is already sorted and it moves $\mathrm{a}[\mathrm{i}]$ to the suitable position.
Program B and C can't sort the number correctly.
Program B is bubble sort as in each i for loop, it move the smallest number in a[i..n 1] into a[i]
12 B There is only two cases to pick two numbers which their product is not a multiple for four.

1. odd x odd
2. $\{2,6\} \times$ odd

So the number of ways is $4 \mathrm{C} 2+2 * 4=14$.
13 A If we place a rook on ( $\mathrm{x}, \mathrm{y}$ ), we cannot place any rook on row x and column y any more. So, we can view place a rook on ( $\mathrm{x}, \mathrm{y}$ ) as delete row x and column y . If the player cannot delete any row or column (all row or column is deleted), he loses.

So we only care about $\min (\mathrm{N}, \mathrm{M})$ as after $\min (\mathrm{N}, \mathrm{M})$ moves, row or column are all deleted and the game is ended. If $\min (N, M)$ is odd, then the first player win, else the first player lose.

Only $\min (3,5)$ is odd, so only when $\mathrm{N}=3$ and $\mathrm{M}=5$, Alice will win.
14 B No matter $\mathrm{f}(\mathrm{x})$ will overflow, $\mathrm{f}(\mathrm{x})$ is always even number.
When $x * 2>=2^{\wedge} 31$, $x$ will overflow and become negative number.
E.g $\mathrm{f}\left(2^{\wedge} 30\right)=2^{\wedge} 31=2147483648=-2147483648$. So (i) is true and (ii) is false.

15 B Notice that we only need to consider the last character since $y$ can always be express in $16 \mathrm{k}+\left(\mathrm{s}[\mathrm{n}-1]-\mathrm{A}^{\prime}\right)$ where 16 k is a even number. So we only care about ( $\mathrm{s}[\mathrm{n}-1]-$ ' $A$ ') is even or not. Only (ii) satisfy it. (' $E$ ' - 'A' $=4=$ even number)
16 C (ii) is not possible because when Alice go, Bob will also go. So there don't exist situation that when Alice go, Bob doesn't go.
17 A Let's rephrase the question in to the following.
Which of the expression is same as ( $x>=y$ )
b) $(-x<=-y)$
c) $(y<x)$
d) ! $(x-y<=y-x)$
$b=(-x<=-y)=(x>=y)$
c) $(y<x)=(x>y)$
$d=!(x-y<=y-x)=!(x<=y)=x>y$

So only $b$ is same as ( $x>=y$ )

18 D Since x and y is unsigned integer. ( $\mathrm{x}||\mid \mathrm{y}$ ) performs logical OR on x 's and y 's bit patterns. So no 1's bit will lost and the result is always $>=x$ and $>=y$.
( $\mathrm{x}||\mid \mathrm{y}$ ) performs logical AND on x's and y's bit patterns. So no 1's bit will gain and the result is always $<=x$ and $<=y$.

So both (i), (ii) and (iii) are true.
19 B Trace the program carefully and will get the following result.
$\mathrm{A}=20$
$B=40$
C $=30$
$\mathrm{D}=30$
20 A The program will output a * $\max (0,(8-\mathrm{b}))$ number of $* \mathrm{~s}$.
So A will output 12 *s. B will output 10 *s. C will output 0 s. D will output 11 *s.
21 D Notice in the j for loop, it compare $\mathrm{a}[\mathrm{i}]$ and $\mathrm{a}[\mathrm{i}+1]$ but not $\mathrm{a}[\mathrm{j}]$ and not $\mathrm{a}[\mathrm{j}+1]$. So it is not bubble sort. Trace the program carefully and you will get the answer.
22 C There is 3 if part in this program.

1) if (score >=90) cout << "A"; else if (score $>=80$ ) cout << "B";
2) if (score $>=70$ ) cout $\ll " \mathrm{C}$ ";
3) if (score >=60) cout << "D"; else cout << "F";

70 satisfy (2) and the first if statement of (3).
So the output is CD.
23 A Notice that $7!=5040>2017$.
So ai is $<7$ for any i from 1 to $n$.

We can use greedy algorithm, which pick the number $k$ from 6 to 0 whenever the current sum $+\mathrm{k}!<=2017$, to find out the solution which n is minimized.

We can get $2 * 6!+4 * 5!+4 * 4!+1!=1440+480+96+1=2017$.

However, since $0!=1!=1$, we can replace $1!$ with $0!$. So the sum of ai $=2 * 6+4 * 5$ $+4 * 4+0=48$
24 C The $3^{\text {rd }}$ line only restrict that data has to be an array and the first element of data is a number. So it CAN store an array of 10 numbers.
25 D a.slice $(\mathrm{x}, \mathrm{y})$ extracts a section of array from $\mathrm{a}[\mathrm{x}]$ to $\mathrm{a}[\mathrm{y}]$. If y is omitted, the a.slice $(\mathrm{x})$ extracts a section of array from $\mathrm{a}[\mathrm{x}]$ to end of array.
So a.slice(0) extract array a from a[0] to its end, which means make a copy of array.

## Section B

| Answer and Explanation |  |  |
| :---: | :---: | :---: |
| A1 | 22 | 62 |
| A2 | $\mathrm{a}[\mathrm{j}+1]:=\mathrm{t}$; | $a[j+1]=t ;$ |
|  | $j$ is the last index where $a[j]<=t$, so we should insert $t$ in $a[j+1]$ |  |
| B | 7 |  |
| C | $753 \mathrm{x}+84 /$ - OR $753 \times 84 /-+$ |  |
| D | 1, 7, 10 (allow permutation) |  |
| E | 3 |  |
| F | st + ed - i |  |
|  | st + ed - i means the last $i^{\text {th }}$ element |  |
| G | true | true or 1 |
|  | A character is always a palindrome |  |
| H | s[st] = s[ed] | s [st] == s[ed] |
|  | A string of two character is palindrome if both characters are equal |  |
| I | ( $s[s t]=s[e d]$ ) and $f(s t+1$, ed - <br> 1) <br> (Accept reverse order) | $s[s t]==s[e d] \& \& f(s t+1$, ed - <br> 1) <br> (Accept reverse order) |
|  | A palindrome is made by two equal character including another palindrome |  |
| J |  |  |
| K | (i + j - 2) mod $a+1$ <br> (May replace a by 5) | $\begin{gathered} (i+j-2) \% \text { a }+1 \\ (\text { May replace a by } 5) \end{gathered}$ |
| L | $x \bmod 10+x$ div $10 \bmod 10(o r x \bmod 100 \operatorname{div} 10)+\mathrm{x}$ div $100=8$ |  |
|  | $x \% 10+x / 10 \% 10$ (or $\mathrm{x} \% 100 / 10)+\mathrm{x} / 100==8$ |  |
|  | $x \% 10$ is the units digit. $x$ div $10 \bmod 10$ is the tens digit. $x$ div 100 is the hundreds digit. |  |
| M | 16 |  |
|  | Count the number of pairs such that $a[j]>a[i]$ and $j<i$ |  |

